

# Sparsity Methods for Systems and Control

## What is Sparsity?

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- 1 Redundant Dictionary
- 2 Underdetermined Systems
- 3 The  $\ell^0$  Norm
- 4 Exhaustive Search

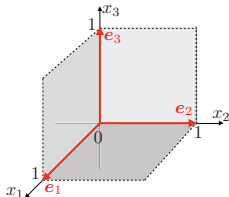
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# Standard basis for $\mathbb{R}^3$

- Standard basis  $\{e_1, e_2, e_3\}$ :

$$e_1 = \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix}, \quad e_2 = \begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix}, \quad e_3 = \begin{bmatrix} 0 \\ 0 \\ 1 \end{bmatrix}.$$



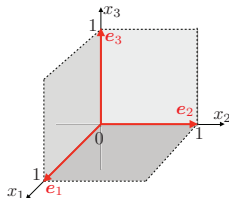
- Any vector  $y \in \mathbb{R}^3$  can be represented as

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# General basis for $\mathbb{R}^3$

- Any three **linearly independent** vectors  $\phi_1$ ,  $\phi_2$ , and  $\phi_3$  in  $\mathbb{R}^3$  form a basis for  $\mathbb{R}^3$ .
- Any vector  $y \in \mathbb{R}^3$  can be represented as

$$y = \beta_1 \phi_1 + \beta_2 \phi_2 + \beta_3 \phi_3$$

- The coefficients  $\beta_1, \beta_2, \beta_3$  are **uniquely determined**.

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# Redundant basis

- Three linearly independent vectors:

$$\phi_1 = e_1 + e_2 = \begin{bmatrix} 1 \\ 1 \\ 0 \end{bmatrix}, \quad \phi_2 = e_2 + e_3 = \begin{bmatrix} 0 \\ 1 \\ 1 \end{bmatrix}, \quad \phi_3 = e_3 + e_1 = \begin{bmatrix} 1 \\ 0 \\ 1 \end{bmatrix}.$$

- Set of 6 vectors (redundant basis)

$$\{e_1, e_2, e_3, \phi_1, \phi_2, \phi_3\}$$

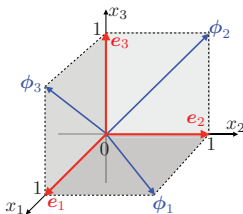
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- Set of 6 vectors (**redundant basis**)

$$\{e_1, e_2, e_3, \phi_1, \phi_2, \phi_3\}$$



- For a vector  $\mathbf{y} \in \mathbb{R}^3$ , we want a signal representation (**redundant representation**):

$$\mathbf{y} = \sum_{i=1}^3 \alpha_i \mathbf{e}_i + \sum_{i=1}^3 \beta_i \phi_i.$$

- There are **infinitely many** solutions for  $\alpha_i$  and  $\beta_i$  ( $i = 1, 2, 3$ ).

$$(\alpha_1, \alpha_2, \alpha_3, \beta_1, \beta_2, \beta_3) = (y_1, y_2, y_3, 0, 0, 0),$$

$$(\alpha_1, \alpha_2, \alpha_3, \beta_1, \beta_2, \beta_3) = (-y_3, -y_1, -y_2, y_1, y_2, y_3).$$

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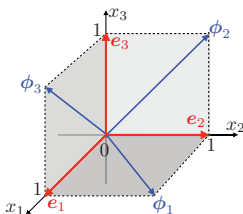
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# Sparse representation

- A vector  $\mathbf{y} = (1, 1, 1)^\top$  on the plane spanned by  $\mathbf{e}_1$  and  $\phi_2$ .
- A coefficient set is obtained as

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This is a **sparse representation** of  $\mathbf{y} = (1, 1, 1)^\top$  since it contains **many zeros**.

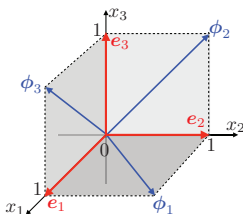


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# Redundant dictionary

- How do you explain this picture by using words in a small dictionary that does not have the word "elephant"?



# Redundant dictionary

- A set of vectors  $\{\phi_1, \phi_2, \dots, \phi_n\}$  in  $\mathbb{R}^m$ .
- If  $m < n$  and  $m$  vectors in this set are linearly independent, then this is called a **redundant dictionary**.
- The elements  $\phi_1, \phi_2, \dots, \phi_n$  in the dictionary is called **atoms** (not “words”).
- For a vector  $y \in \mathbb{R}^m$ , we find coefficients  $\alpha_1, \alpha_2, \dots, \alpha_n$  such that

$$y = \sum_{i=1}^n \alpha_i \phi_i.$$

- If the dictionary is more redundant, then we may obtain a sparser coefficient set.



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# Sparse representation

- Define a matrix  $\Phi$  and a vector  $\mathbf{x}$  as

$$\Phi \triangleq [\phi_1 \quad \phi_2 \quad \dots \quad \phi_n] \in \mathbb{R}^{m \times n}, \quad \mathbf{x} \triangleq \begin{bmatrix} \alpha_1 \\ \alpha_2 \\ \vdots \\ \alpha_n \end{bmatrix} \in \mathbb{R}^n.$$

- Then the relation

$$\mathbf{y} = \sum_{i=1}^n \alpha_i \phi_i.$$

is compactly rewritten as

$$\mathbf{y} = \Phi \mathbf{x}.$$

- The matrix  $\Phi$  is called a **dictionary matrix** or **measurement matrix**.

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# The problem of sparse representation

## Problem (Sparse Representation)

Given a vector  $\mathbf{y} \in \mathbb{R}^m$  and a dictionary matrix  $\Phi \in \mathbb{R}^{m \times n}$  with  $m < n$ . Find the simplest (i.e. sparsest) representation of  $\mathbf{y}$  that satisfies

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# Linear equations

- linear equations with unknowns  $x_1$ ,  $x_2$ , and  $x_3$ :

$$x_1 + x_2 + x_3 = 3$$

$$x_1 - x_3 = 0$$

- There are infinitely many solutions
- All solutions

$$x_1 = t, \quad x_2 = -2t + 3, \quad x_3 = t,$$

where  $t \in \mathbb{R}$  is a parameter.

- Such a system of equations is called an **underdetermined system**.
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## Minimum solution

- Let us consider a detective, like Edogawa Conan<sup>1</sup>, who solve this problem.
- The two proofs (equations) are insufficient and he should seek one more **independent** proof.
- If he gets one more proof saying *the criminal is the smallest one among the suspects*.
- The  $\ell^2$ -norm

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$$\begin{aligned}\|x\|_2^2 &= x_1^2 + x_2^2 + x_3^2 \\ &= t^2 + (-2t + 3)^2 + t^2 \\ &= 6(t - 1)^2 + 3.\end{aligned}$$

is minimized by  $t = 1$ .

- The unique solution is  $(x_1, x_2, x_3) = (1, 1, 1)$ .

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# Linear equations in matrix form

- Linear equations in a matrix form:

$$\Phi \mathbf{x} = \mathbf{y}.$$

- $\Phi$  is an  $m \times n$  matrix where  $m < n$  (we call this a **fat matrix**).
- Assume  $\Phi$  has **full row rank**, that is,

$$\text{rank}(\Phi) = m.$$

- For any vector  $\mathbf{y} \in \mathbb{R}^m$ , there exists at least one solution  $\mathbf{x}$  that satisfies  $\Phi \mathbf{x} = \mathbf{y}$ .
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## Norm

A norm in  $\mathbb{R}^n$  should satisfy

- 1 For any vector  $\mathbf{x} \in \mathbb{R}^n$  and any number  $\alpha \in \mathbb{R}$ ,  $\|\alpha\mathbf{x}\| = |\alpha|\|\mathbf{x}\|$ .
- 2 For any  $\mathbf{x}, \mathbf{y} \in \mathbb{R}^n$ ,  $\|\mathbf{x} + \mathbf{y}\| \leq \|\mathbf{x}\| + \|\mathbf{y}\|$ .
- 3  $\|\mathbf{x}\| = 0 \iff \mathbf{x} = \mathbf{0}$ .

- The  $\ell^2$  norm (or Euclidean norm)

$$\|\mathbf{x}\|_2 \triangleq \sqrt{x_1^2 + x_2^2 + \cdots + x_n^2}.$$

- The  $\ell^1$  norm

$$\|\mathbf{x}\|_1 \triangleq |x_1| + |x_2| + \cdots + |x_n|.$$

- The  $\ell^\infty$  norm (or maximum norm)

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- 2 For any  $x, y \in \mathbb{R}^n$ ,  $\|x + y\| \leq \|x\| + \|y\|$ .
- 3  $\|x\| = 0 \iff x = 0$ .

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$$\|x\|_2 \triangleq \sqrt{x_1^2 + x_2^2 + \cdots + x_n^2}.$$

- The  $\ell^1$  norm

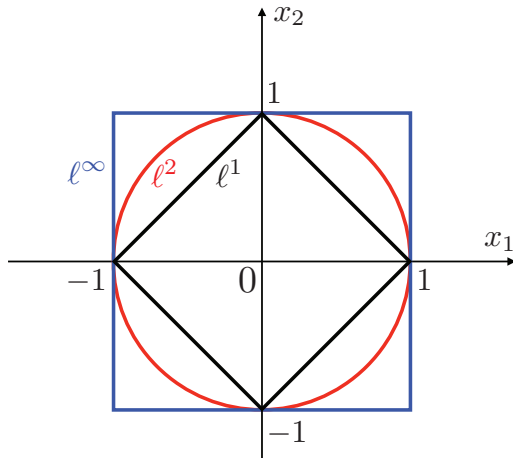
$$\|x\|_1 \triangleq |x_1| + |x_2| + \cdots + |x_n|.$$

- The  $\ell^\infty$  norm (or maximum norm)

$$\|x\|_\infty \triangleq \max\{|x_1|, |x_2|, \dots, |x_n|\}$$

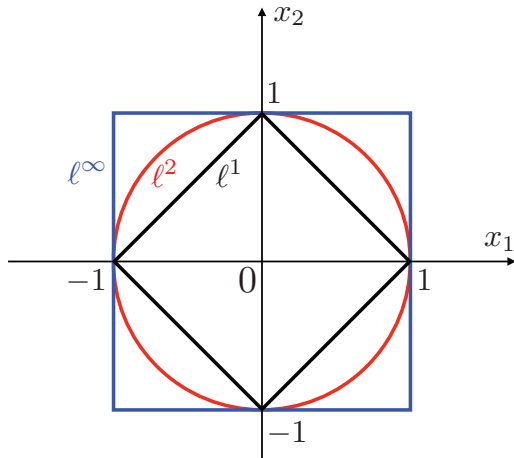
# Norms in $\mathbb{R}^n$

Contour curves ( $\|x\|_p = 1$ ) of  $\ell^1$ ,  $\ell^2$ ,  $\ell^\infty$  norms.



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# $\ell^0$ norm

- Consider a vector  $x = [x_1, x_2, \dots, x_n]^T \in \mathbb{R}^n$ .
- The  $\ell^0$  norm of  $x$  is defined by

$$\|x\|_0 \triangleq |\text{supp}(x)|,$$

where

- $\text{supp}$  is the support of  $x$ , namely,

$$\text{supp}(x) \triangleq \{i \in \{1, 2, \dots, n\} : x_i \neq 0\},$$

- $|\cdot|$  denotes the number of elements.

- The  $\ell^0$  norm counts the number of non-zero elements in  $x$ .
- The  $\ell^0$  norm *does not* satisfy the first property in the definition of norm, and it is sometimes called the  $\ell^0$  pseudo-norm.

$$\|2x\|_0 = \|x\|_0 \neq 2\|x\|_0.$$

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# $\ell^0$ optimization problem

Now the problem of sparse representation is formulated as follows:

## $\ell^0$ optimization problem

Given a vector  $\mathbf{y} \in \mathbb{R}^m$  and a full-row-rank matrix  $\Phi \in \mathbb{R}^{m \times n}$  with  $m < n$ . Find the optimizer  $\mathbf{x}^*$  of the optimization problem:

$$\underset{\mathbf{x} \in \mathbb{R}^n}{\text{minimize}} \quad \|\mathbf{x}\|_0 \quad \text{subject to} \quad \mathbf{y} = \Phi \mathbf{x}.$$

This problem is called the  $\ell^0$  optimization.

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- 1 Redundant Dictionary
- 2 Underdetermined Systems
- 3 The  $\ell^0$  Norm
- 4 Exhaustive Search



# How to solve it?

## $\ell^0$ optimization problem

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- We can try an **exhaustive search** for this optimization.

# Exhaustive search: example

- Find the minimum- $\ell^0$  solution  $(x_1, x_2, x_3)$  that satisfies

$$x_1 + x_2 + x_3 = 3$$

$$x_1 - x_3 = 0$$

- First, try  $(x_1, x_2, x_3) = (0, 0, 0)$ . This is not a solution.
- Second, try  $(x_1, 0, 0)$ ,  $(0, x_2, 0)$ , and  $(0, 0, x_3)$ .

• If  $(x_1, 0, 0)$  is a solution, then  $x_1 = 3$  and  $x_3 = 0$ . This is not a solution.

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# Exhaustive Search (step 1)

- If  $y = 0$ , then output  $x^* = 0$  as the optimal solution and quit.
- Otherwise, proceed to the next step.

# Exhaustive Search (step 2)

- Find a vector  $\mathbf{x}$  with  $\|\mathbf{x}\|_0 = 1$  that satisfies the equation  $\mathbf{y} = \Phi\mathbf{x}$ .

That is, set

$$\mathbf{x}_1 \triangleq \begin{bmatrix} x_1 \\ 0 \\ \vdots \\ 0 \end{bmatrix}, \quad \mathbf{x}_2 \triangleq \begin{bmatrix} 0 \\ x_2 \\ 0 \\ \vdots \\ 0 \end{bmatrix}, \quad \dots, \quad \mathbf{x}_n \triangleq \begin{bmatrix} 0 \\ \vdots \\ 0 \\ x_n \end{bmatrix}$$

and search  $x_i \in \mathbb{R}$  ( $i = 1, 2, \dots, n$ ) that satisfies

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If a solution exists for any  $i$ , output  $\mathbf{x}^* = \mathbf{x}_i$  as the solution and quit.

Otherwise, output  $\mathbf{x}^* = \mathbf{0}$  as the solution.

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- Otherwise, proceed the next step.

# Exhaustive Search (step $k$ )

- Do similar procedures for  $\|x\|_0 = k, k = 3, 4, \dots, m$ .

# Is exhaustive search useful?

- It is easily implemented.
- The computation time to find a solution grows **exponentially** with problem size  $m$ .
- Suppose  $m = 100$ .
  - Then it roughly takes  $2^{100} = 1.3 \times 10^{30}$  iterations (at the worst).
  - If we can do one iteration in  $10^{-7}$  seconds (by a super computer),
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- The computation time to find a solution grows **exponentially** with problem size  $m$ .
- Suppose  $m = 100$ .
  - Then it roughly takes  $2^{100} \approx 1.3 \times 10^{30}$  iterations (at the worst).
  - If we can do one iteration in  $10^{-15}$  seconds (by a super computer),
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# Conclusion

- Sparsity of a vector is measured by its  $\ell^0$  norm.
- In sparse representation, a redundant dictionary of vectors is used.
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